

Network Modeling and Barabasi-Albert (BA) Model

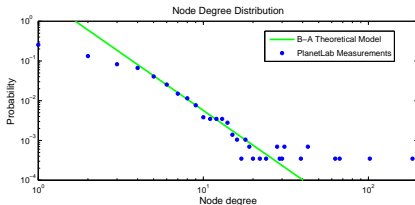
- ▶ *Scale free* model for simulation of node degree distribution

- ▶ growth
- ▶ preferential attachment

$$P(k) = 2m^{1/\beta} k^{-\gamma} \text{ for } \gamma=3, \beta=0.5$$

- ▶ *Clustering* coefficient

$$C_i = \frac{2E_i}{k_i(k_i-1)}$$

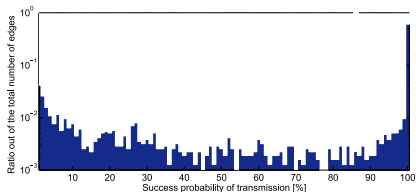


Node degree distribution - B-A Model vs PlanetLab



Real Scale Free Networks - PlanetLab Measurements

- ▶ Internet testbed – 226 end hosts
- ▶ Measured topology and packet loss information (*ping*, *traceroute*).
- ▶ IP aliasing resolution

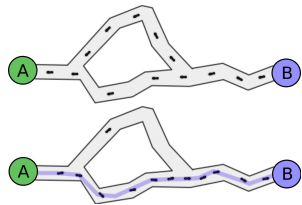


Derived link loss (*erasure*) distribution based on path loss measurements, employing network tomography techniques (Netscope).



Ant Colony Optimization for Finding Shortest Path

- ▶ Ants walk through the network, looking for food sources. Once a food source is found, ants return to the colony, on the same path, laying down a *pheromone* trail.
- ▶ Each ant has a memory of the current path.
- ▶ Once food is found, ants return to colony while laying down a pheromone trail.
- ▶ At every node, an ant chooses one among the adjacent links based on previous pheromone levels of the links.



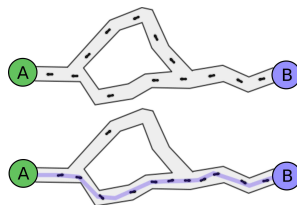
Ant Colony Optimization for Finding Shortest Path

- ▶ Transition probability

$$p_{ij}(t) = \frac{\tau_{ij}^{\alpha}(t)\eta_{ij}^{\beta}}{\sum_{s \text{ allowed}} \tau_{is}^{\alpha}(t)\eta_{is}^{\beta}}$$

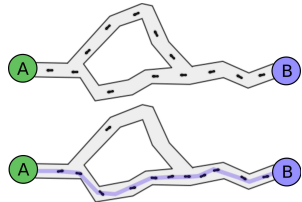
- ▶ Pheromone update

$$\tau_{ij}(t+n) = \rho\tau_{ij}(t) + \Delta\tau_{ij},$$



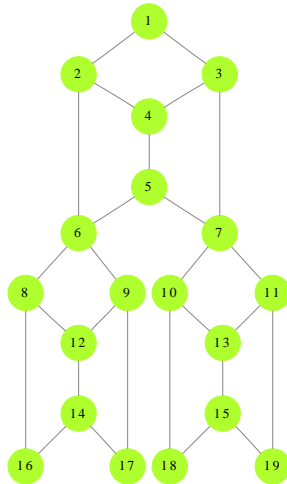
Ant Colony Optimization for Finding Shortest Path

- ▶ Since shorter paths will be walked more times in the same interval, a higher pheromone concentration is noticed.
- ▶ Ant Colony Optimization (ACO) is preferable among other shortest path algorithms since nodes do not need to be aware of the network topology.



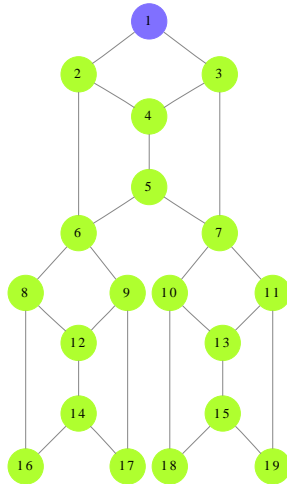
Finding Disjoint Paths

- ▶ Select a source node and sink nodes.



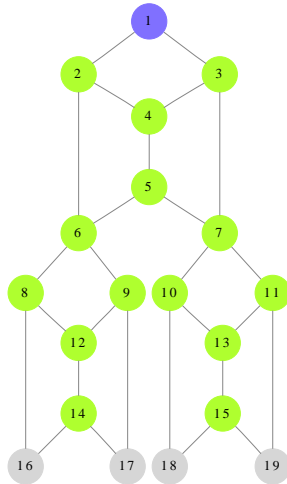
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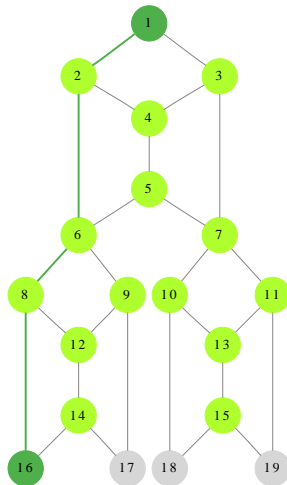
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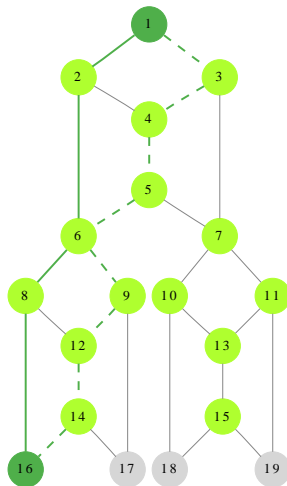
Finding Disjoint Paths

- ▶ Get multicast rate m_i for every pair (*source*, *sink_i*). Overall multicast rate is $\min(m_i)$.
- ▶ For all sinks, find $\min(m_i)$ disjoint paths.



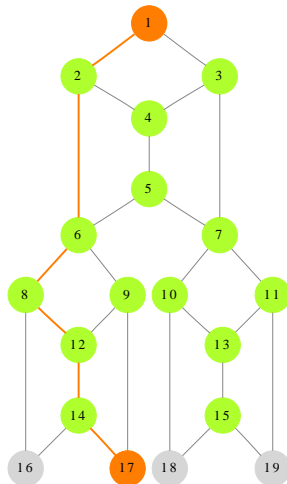
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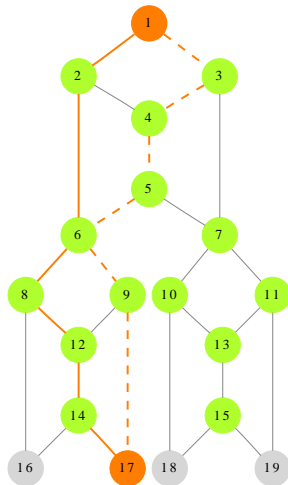
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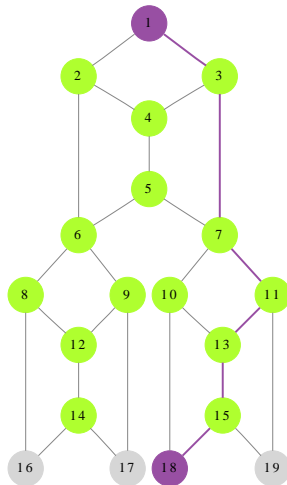
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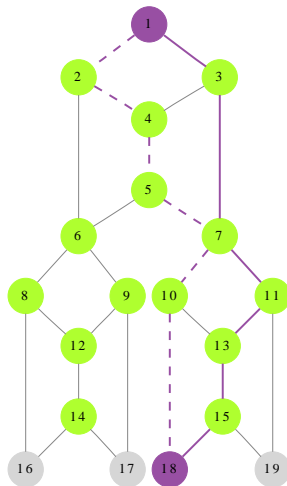
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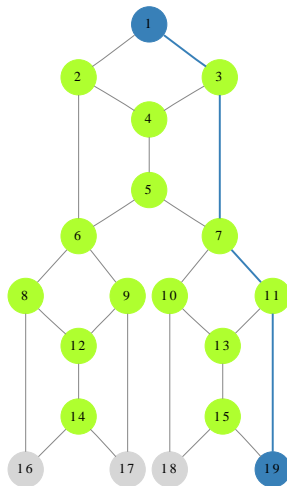
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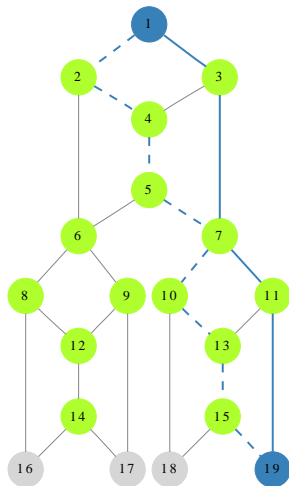
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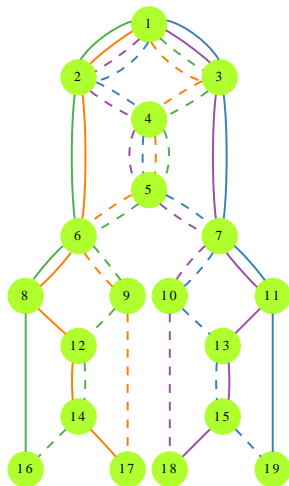
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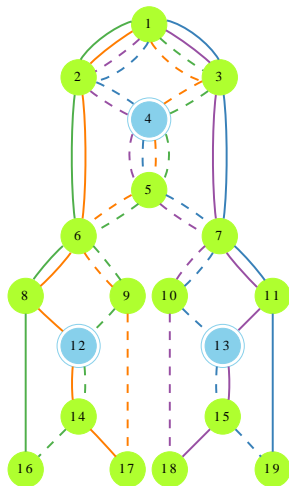
Finding Coding Nodes

- ▶ Obtain coding nodes from the intersection of paths corresponding to different sinks.
- ▶ Ensure that at all times each edge passes either coded or not coded messages, not both in the same time, otherwise flip disjoint paths.



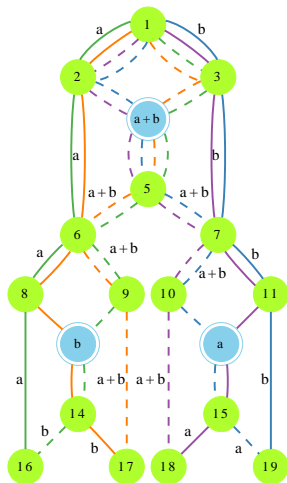
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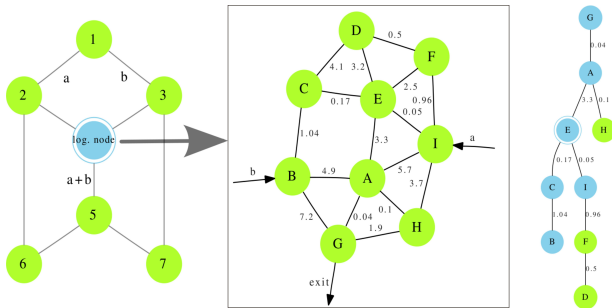


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Hierarchical Network Structure



► *Logical node*

► Minimum Spanning Tree

